

# CSI 33 LECTURE NOTES (Ojakian)

## Topic 8: Heaps

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### OUTLINE

(References: 13.1, 13.2)

1. Heaps
  2. Priority Queue
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### 1. Priority Queue

- (a) What it is: See interface on page 445.

Note: We will ignore the “item” in enqueue.

**PROBLEM 1.** *Determine the result of the dequeues:*

```
enqueue(5)
enqueue(2)
enqueue(3)
enqueue(9)
```

```
dequeue()
dequeue()
dequeue()
dequeue()
```

- (b) Examples of application: See page 444: Hospital Triage and Computer Operating System Tasks.

- (c) Problems implementing it as a list:

- i. If Enqueue at back, then Dequeue is inefficient
- ii. If Enqueue in order, then Enqueue inefficient (finding spot could be fast, but may need to shift)

**PROBLEM 2.** *Observe the above two inefficient approaches on the above example.*

- (d) What about Binary Search Tree?

Only good as long as balanced...

**PROBLEM 3.** *Consider the operation of a priority queue using a BST in the following example.*

```
enqueue(2)
enqueue(3)
enqueue(5)
enqueue(9)
```

```
dequeue()
dequeue()
dequeue()
dequeue()
```

(e) A response: Heaps! - key point: it will be forced to remain balanced.

## 2. Heap

Note: Also called Binary Heap or Max Heap

(a) Definition (Do via example on page 446):

- i. Binary Tree with ordered labels
- ii. It's complete
- iii. The value at any node is as large or larger than its children
- iv. Note: This is a **MAX heap**. Also exists **MIN heap**
- v.

**PROBLEM 4.** *Draw examples. Which are Max Heaps and which are Min Heaps and which are neither?*

(b) Insertion (do via example on page 448):

First consider a problematic approach: Start at the top and move down to insertion spot similar to BST What could go wrong?

- i. Place the item at the bottom in the “next” spot
- ii. Swap it up the heap till it is correctly positioned

(c) Pop Max (do via example on page 447):

First consider a problematic approach: Successively move larger child up the tree to fill the deleted spot. What could go wrong?

- i. Remove the “last” node in the tree.
- ii. Replace the root (i.e. the max) with this node.
- iii. Swap it down the heap till it is correctly positioned

### 3. Implementation

- (a) Difficulty in using our trees based on link structures: How do you deal with the “last” node.

**PROBLEM 5.** Consider how you would find the last node in a linked-based binary tree.

**PROBLEM 6.** Suppose you tried to correct, by maintaining a pointer to the last node. Observe how the last node can change in somewhat complicated ways.

- (b) Alternative- Binary trees represented as an array (See Chapter 7, pages 229-230).

**PROBLEM 7.** Draw a (non-complete) binary tree and give its array representation in two ways: starting at index 0 and starting at index 1.

**PROBLEM 8.** Based on the last example, determine formulas (in both cases) for finding the following: parent, left child, right child.

**PROBLEM 9.** Consider the following array: [NONE, 3, 4, 5, 7, 6, 7, 9, 8, 8]

i. Draw the binary tree represented by this array representation.

ii. What kind of data type have we drawn?

**PROBLEM 10.** Consider the BST we get by the following sequence of inserts: 1, then 2, then 3, then 4.

i. Draw the resulting BST

ii. Using a linked binary tree, how many Tree Nodes of memory are used?

iii. Draw the array representation of this tree.

**PROBLEM 11.** Compare the advantages and disadvantages of representing a binary tree as an array versus a linked structure?

**PROBLEM 12.** Program the binary tree using an array with a start index of 1 (for HW you will modify it to do a start index of 0).

- (c) **Program the Heap with insert and pop max!**