

## 11.1 Trees

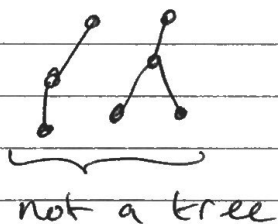
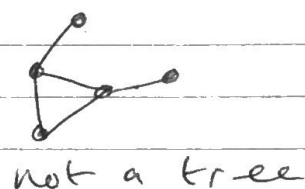
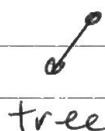
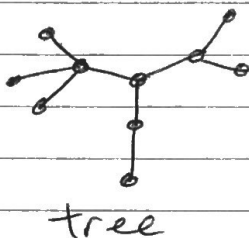
①

Trees are a special type of graph.

Recall the definitions for path, circuit, simple circuit and connected.

Definition: A tree is a connected graph with no simple circuits.

Examples



← this is a collection of 2 trees

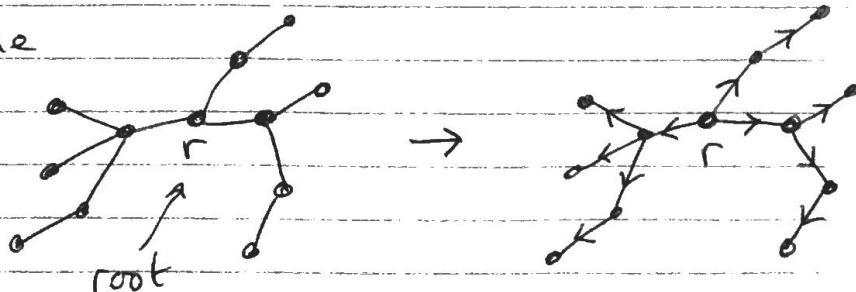
(Collections of trees are called forests).

Property of trees: There is a unique simple path between any two vertices.

Definition: A rooted tree is just a tree where one vertex has been labelled the root.

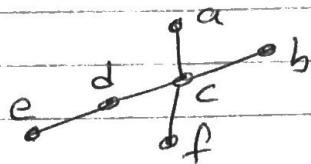
This makes the tree into a directed graph since we can direct every edge away from the root.

For example

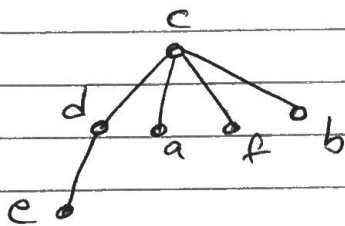


If you choose a different vertex for the root you will get a different directed graph. Usually we draw it so that the root is at the top and all arrows point down.

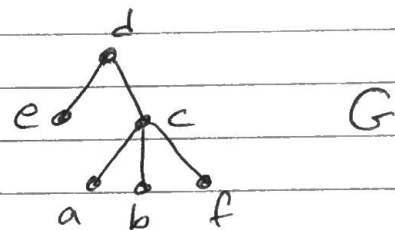
Example



choose c as root  
get



choose d as root



For a rooted tree such as  $G$  there is a lot of terminology to describe the vertices. See pages 747-748 of the text for

- parent
- child
- siblings
- ancestors
- descendants
- leaf
- internal vertex
- subtree rooted at a vertex

} think of a family tree.

For example, in  $G$ ,  $c$  is the parent of  $a, b, f$  and  $a, b, f$  are the children of  $c$ . Also  $c$  and  $d$  are the ancestors of  $a$ .

The leaves of  $G$  are  $e, a, b, f$  and the internal vertices are  $c, d$ .

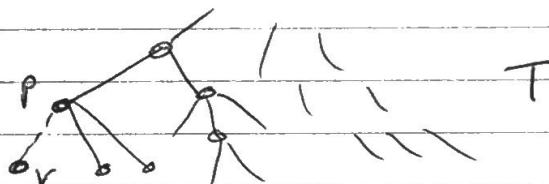
(2)

Theorem A tree with  $n \geq 1$  vertices has  $n-1$  edges.

Proof. Let  $P(n)$  be this statement and we'll prove it by induction.

Basis case:  $P(1)$  is true tree = •

Inductive step: Assume true  $P(k)$  and use to prove  $P(k+1)$ . By choosing a root we may assume a tree  $T$  with  $k+1$  vertices is rooted. It must have a leaf  $v$  with parent  $p$ .

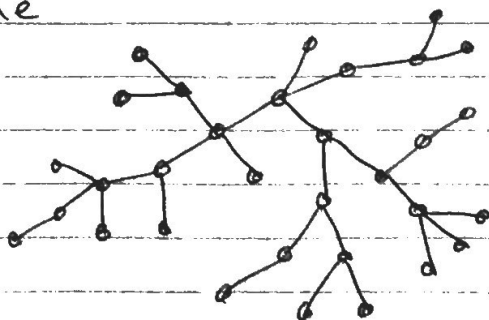


Remove the edge  $(p, v)$  and the vertex  $v$ . A tree with  $k$  vertices remains, so by  $P(k)$  it has  $k-1$  edges. Therefore  $T$  has  $k$  edges. This proves the induction step.

So  $P(n)$  is true for all  $n \geq 1$ .  $\square$

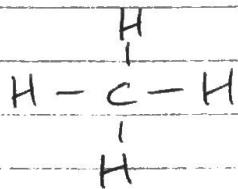
The converse of this theorem is also true: A connected graph with  $n$  vertices and  $n-1$  edges must be a tree.

Example

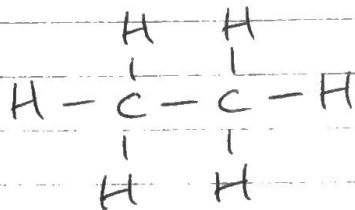


vertices?  
edges?

## Hydrocarbon models



methane



ethane

These are examples of saturated hydrocarbon molecules  $\text{C}_n \text{H}_{2n+2}$  (see p 750).

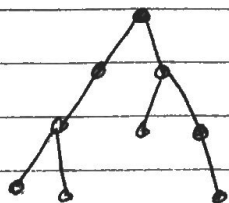
Each carbon is a vertex with degree 4 and each Hydrogen a vertex with degree 1.

So  $\text{C}_n \text{H}_{2n+2}$  has  $3n+2$  vertices. Its total degree is  $4n + 1(2n+2)$ . This is twice the number of edges by the Handshaking theorem, so  $3n+1$  edges.

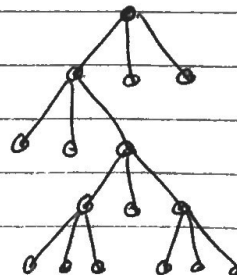
By our previous result  $\text{C}_n \text{H}_{2n+2}$  must be a tree.

Definition: A rooted tree is called an m-ary tree if every internal vertex has at most  $m$  children. If they have exactly  $m$  children then called a full m-ary tree.

examples



2-ary tree  
(binary tree)



a full 3-ary  
tree



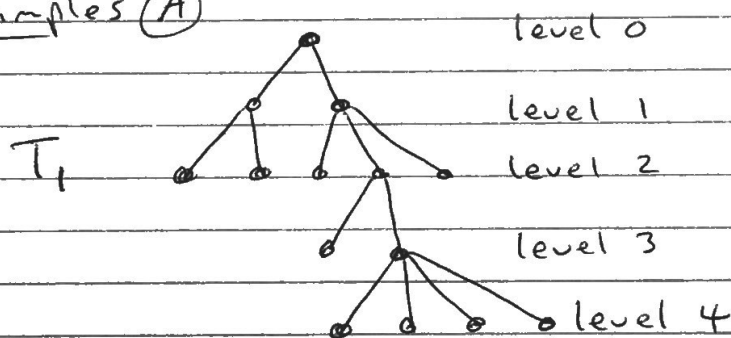
Some last terminology for rooted trees.

The level of a vertex is the length of the path from that vertex to the root.

The height of a rooted tree is the biggest level of any vertex.

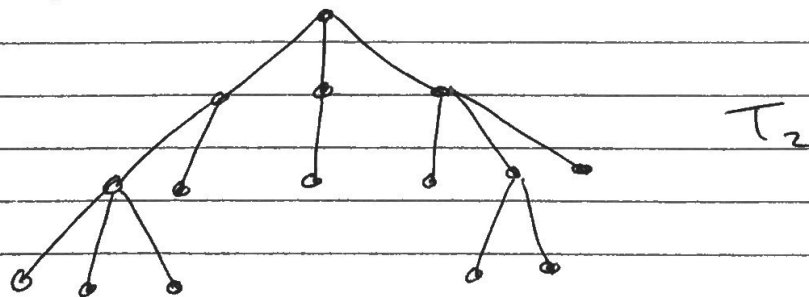
Suppose a rooted tree has height  $h$ . Then it is balanced if all leaves are at level  $h$  or  $h-1$ .

Examples (A)



This tree  $T_1$  has height 4 and is not balanced.

Example (B)



This tree  $T_2$  has height 3 and is balanced.  
It is a 3-ary tree.