

Review of chapter 1.

1.1, 1.2 Propositional Logic

A proposition p must be true (T) or false (F).

We can combine propositions using the logical operations:

\wedge	and
\vee	or
\oplus	exclusive or (xor)
\rightarrow	implication
\leftrightarrow	if and only if
\neg	not

Example ① Use propositions and logical operations to express this "I will go for a walk if it is warm and not raining."

Solution: We can use these propositions
 p says "I will go for a walk"
 q says "It is warm"
 r says "It is raining".

Clearer to rewrite the sentence as

"If it is warm and it is not raining then I will go for a walk"

So we want $(q \wedge \neg r) \rightarrow p$

Example (2) Suppose p is true and q is false. Find the truth value of $\neg(p \rightarrow q) \vee q$

Solution: First $p \rightarrow q$ is false so $\neg(p \rightarrow q)$ is true. Then $T \vee F$ is true. So $\neg(p \rightarrow q) \vee q$ is true.

An implication $p \rightarrow q$ has three related statements:

$q \rightarrow p$	converse
$\neg q \rightarrow \neg p$	contrapositive
$\neg p \rightarrow \neg q$	inverse

but only the contrapositive is logically equivalent to the first implication.

Example (3) Give the inverse of this statement "If it is not Sunday then the post office is open."

Solution: Inverse is "If it is Sunday then the post office is closed."

On an island everyone is either a knight who only says true things or a knave who only says false things. Use logic to solve the next example.

2.

Example (4) You meet two people on the island. A says "I am a knight and B is a knave". B says "we're both knights". Decide if A and B are knights or knaves.

Solution: We see they both cannot be telling the truth, so one at least is a knave. Therefore B must be a knave. A can be a knight or a knave - both cases are consistent with what they say.

1.3 Propositional equivalences

Combinations of propositions that are always true are called tautologies.

Example (5) Show that $(p \wedge \neg p) \rightarrow p$ is a tautology.

Solution: First look at $p \wedge \neg p$. If p is T or F we have $p \wedge \neg p$ being false. So $(p \wedge \neg p) \rightarrow p$ becomes

$$F \rightarrow p$$

and this is always true (by the definition of implication). So $(p \wedge \neg p) \rightarrow p$ is a tautology.

Two propositions p, q are logically equivalent if they always have the same truth value.

We write $p \equiv q$

if $p \leftrightarrow q$ is a tautology.

Example (6) Use a truth table to show that $\neg(p \wedge q) \equiv \neg p \vee \neg q$

Solution

p	q	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
T	T	T	F	F	F	F
T	F	F	T	F	T	T
F	T	F	T	T	F	T
F	F	F	T	T	T	T

these are the same means we have shown the equivalence

$\neg(p \wedge q) \leftrightarrow \neg p \vee \neg q$
is always true.

Some laws of logic

- $\neg(\neg p) \equiv p$
 - $p \wedge q \equiv q \wedge p$
 - $p \wedge (q \wedge r) \equiv (p \wedge q) \wedge r$
 - $\neg(p \wedge q) \equiv \neg p \vee \neg q$
 - $\neg(p \vee q) \equiv \neg p \wedge \neg q$
- } De Morgan
- $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$

1.4, 1.5 Quantifiers

Propositional functions $P(x)$ are true or false depending on what x is. The set of possible x s is called the domain.

Example (7) let the domain be all real numbers and let $P(x)$ say " $x^2 + 1 = 10$ ". Is $P(10)$ true? Can we find an x that makes $P(x)$ true?

Solution: $P(10)$ says $10^2 + 1 = 10$
 $101 = 10$

and that is false. We can find a real number that works: $P(3)$ is true.

$\forall x P(x)$ says "for all x in the domain we have $P(x)$ "

and for example (7) it means

"for all real numbers x we have $x^2 + 1 = 10$ ".

That's false.

$\exists x P(x)$ says "there exists an x in the domain so that $P(x)$ ".

and for example (7) it means

"there exists a real number x so that $x^2 + 1 = 10$ ".

That's true.

We can have more than one variable and nested quantifiers:

$$\forall x \exists y (3x = y) \quad \text{domain} = \text{reals}$$

means "for all real numbers x there exists a real y so that $3x = y$ ".

This is true - starting with any x we can find y .

There is a way to negate quantifiers. The negation of the last statement is

$$\begin{aligned} \neg (\forall x \exists y (3x = y)) \\ \equiv \exists x \neg (\exists y (3x = y)) \\ \equiv \exists x \forall y \underbrace{\neg (3x = y)}_{3x \neq y} \end{aligned}$$

negation is

"there exists a real x so that for all real y we have $3x$ not equal to y "

This is false.

1.6 Rules of inference

Review modus ponens, modus tollens, hypothetical syllogism and the other rules of inference in this section.