

CSI33 DATA STRUCTURES

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OUTLINE

- 1 C++ SUPPLEMENT: 1.2
 - A Binary Search Tree



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 - A Binary Search Tree

THE BINARY SEARCH PROPERTY

A BINARY TREE IS “SORTED”

A **Binary Search Tree**, or BST, is a binary tree where every node has the following property:

- Each value in the left subtree is less than the value at the node.
- Each value in the right subtree is greater than the value at the node.



THE BINARY SEARCH PROPERTY

BINARY SEARCH WITH A BINARY TREE

- Start at the root
- If the value is there, you are done
- If the value is less than the node value, search the left subtree
- If the value is greater than the node value, search the right subtree



THE BINARY SEARCH PROPERTY

PERFORMANCE (RUNNING TIME) TO FIND A VALUE

- Average Performance Is $\Theta(\log n)$. If the tree is not too unbalanced, then we divide the number of items to search in half at each node. This is actually a binary search.
- Worst-Case Performance Is $\Theta(n)$. If the tree branches only to one side (left or right) this is the same as linear search.



IMPLEMENTING A BST

BST.H (HEADER FILE) I

```
class BST
{
public:
    BST();
    ~BST();
    void insert(int value);
    int find(int value);
    void delete_(int value);
```



IMPLEMENTING A BST

BST.H (HEADER FILE) II

```
private:
    TreeNode *_root;
    void _deleteNodes(TreeNode *t);
    TreeNode *_insertRec(TreeNode* t, int value);
    TreeNode *_findRec(TreeNode* t, int value);
    TreeNode *_deleteRec(TreeNode* t, int value);
    TreeNode *_deleteMax(TreeNode *pNode, int& item);
        // updates item
}
```



IMPLEMENTING A BST

BST.CPP (CONSTRUCTOR)

```
BST::BST()
{
    /*
    post:  creates empty binary search tree
    */
    _root = NULL;
}
```



IMPLEMENTING A BST

BST.CPP (INSERT)

```
void BST::insert(int value)
{
    _root = _insertRec(_root, value);
}
TreeNode *BST::_insertRec(TreeNode* t, int value)
{
    if (t == NULL)
        t = new TreeNode(value, NULL, NULL);
    else if (value < t->_item)
        t->_left = _insertRec(t->_left, value);
    else if (value > t->_item)
        t->_right = _insertRec(t->_right, value);
    else
        // raise exception since value is already in tree
        return t;
}
```



IMPLEMENTING A BST

BST.CPP (FIND)

```
int BST::find(int value)
{
    TreeNode *t = _findRec(_root, value);
    return t->_item;
}
```



IMPLEMENTING A BST

BST.CPP (_FINDREC)

```
TreeNode *BST::_findRec(TreeNode *t, int value)
{
    if (t == NULL) //raise exception since value not in tree
    else
    {
        if (value < t->_item)
            return _findRec(t->_left, value);
        else if (value > t->_item)
            return _findRec(t->_right, value);
        else // found value
            return t;
    }
}
```



IMPLEMENTING A BST

BST.CPP (DELETE_)

```
void BST::delete_(int value)
{
    /*
    pre:  item is in tree
    post: item is not in tree
    */
    _root = _deleteRec(_root, value);
}
```



IMPLEMENTING A BST

BST.CPP (_DELETEREC I)

```
TreeNode *BST::_deleteRec(TreeNode *t, int value)
{
    if (t == NULL)
        //raise exception since value not in tree
    else if (value < t->_item) // deletion from left
        t->_left = _deleteRec(t->_left, value);
    else if (value > t->_item) // deletion from right
        t->_right = _deleteRec(t->_right, value);
    else // delete t itself
```



IMPLEMENTING A BST

BST.CPP (_DELETEREC II)

```
else // delete t itself
    if (t->_left == NULL) //promote right subtree
    {
        TreeNode *new_t = t->_right;
        delete t;
        t = new_t;
    }
    else if (t->_right == NULL) //promote left subtree
    {
        TreeNode *new_t = t->_left;
        delete t;
        t = new_t;
    }
    else // can't delete t--overwrite with max left value
        t->_left = _deleteMax(t->_left, t->_item);
return t;
}
```



IMPLEMENTING A BST

BST.CPP (_DELETEMAX)

```
// use pass-by-reference to overwrite deleted item value
TreeNode *BST::_deleteMax(TreeNode *t, int& item)
{
    if (t->_right == NULL) // t is the max
    {
        TreeNode *new_t = t->_left;
        item = t->_item;
        delete t;
        return new_t;
    }
    else // max in rt subtree--recursively find and delete it
    {
        t->_right = _deleteMax(t->_right, item);
        return t;
    }
}
```



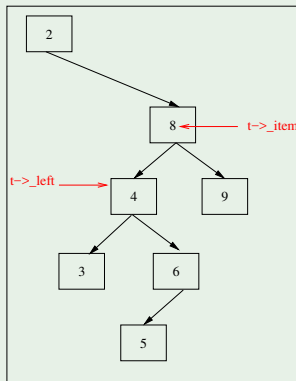
DELETION EXAMPLE: DELETE_(8)

DELETION: NEITHER CHILD IS NULL

```

else // 'delete' t
{
    if (t->_left == NULL)
    { //promote right subtree
        TreeNode *new_t = t->_right;
        delete t;
        t = new_t;
    }
    else if (t->_right == NULL)
    { //promote left subtree
        TreeNode *new_t = t->_left;
        delete t;
        t = new_t;
    }
    else // can't delete t-overwrite max
        t->_left = _deleteMax(t->_left,
            t->_item);
}
return t;

```



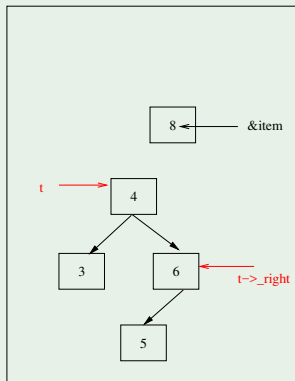
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DELETION: NEITHER CHILD IS NULL

```

TreeNode *BST::_deleteMax(TreeNode *t,
int& item)
{
    if (t->_right == NULL) // t has the max
    {
        TreeNode *max_left = t->_left;
        item = t->_item;
        delete t;
        return max_left;
    }
    else // max is in right subtree
    {
        t->_right = _deleteMax(t->_right,
            item);
        return t;
    }
}

```



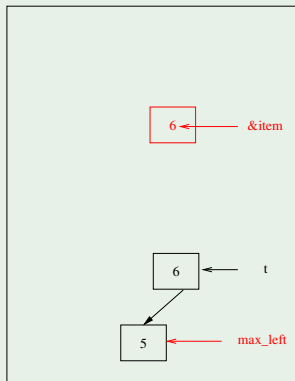
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        item = t->_item;
        delete t;
        return max_left;
    }
    else // max is in right subtree
    {
        t->_right = _deleteMax(t->_right,
            item);
        return t;
    }
}

```



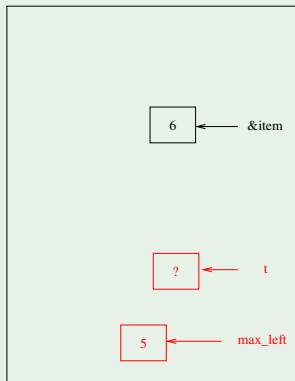
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        delete t;
        return max_left;
    }
    else // max is in right subtree
    {
        t->_right = _deleteMax(t->_right,
            item);
        return t;
    }
}

```



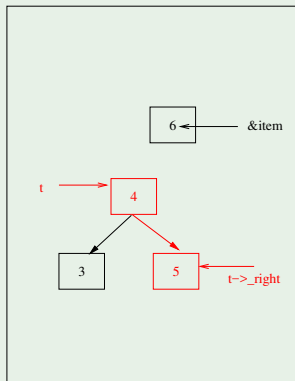
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int& item)
{
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    {
        TreeNode *max_left = t->_left;
        item = t->_item;
        delete t;
        return max_left;
    }
    else // max is in right subtree
    {
        t->_right = _deleteMax(t->_right,
            item);
        return t;
    }
}

```



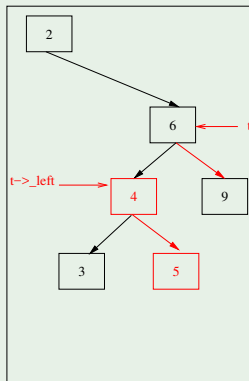
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DELETION: NEITHER CHILD IS NULL

```

else // 'delete' t
{
    if (t->_left == NULL)
    { //promote right subtree
        TreeNode *new_t = t->_right;
        delete t;
        t = new_t;
    }
    else if (t->_right == NULL)
    { //promote left subtree
        TreeNode *new_t = t->_left;
        delete t;
        t = new_t;
    }
    else // can't delete t-overwrite max
        t->_left = _deleteMax(t->_left,
            t->_item);
}
return t;

```



TRAVERSING A BST

COPY DATA INTO AN ARRAY

- Use inorder traversal to keep items in order.
- Then process the array.
- Disadvantage: uses extra memory for the array.



RUN-TIME ANALYSIS OF BST METHODS

METHODS

- **visit** is $\Theta(n)$.
- **insert, delete, find** have $\Theta(\log n)$ average behavior.
- **insert, delete, find** have $\Theta(n)$ worst-case behavior.

